

KATEGORIES & MORFISMS (summary)

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1. INTRODUCTION

The intent is to model the properties of categories and morphisms using *Kategorien* and *morfisms*.

2. KATEGORIES

Kategorien relate properties of *objects* using *morfisms*. A specific *category* is defined to be a collection of certain objects and morfisms following prescribed conditions and rules.

Definition 2.1. Category:

A collection, \mathbf{Q} , of *objects* and *morfisms* qualifies as a *Category* if and only if the following conditions and rules hold.

CONDITIONS:

- (1) There exists a collection of Objects, $\mathbf{Obj}(\mathbf{Q})$, denoted by $\mathbf{A}, \mathbf{B}, \mathbf{C}, \dots$
- (2) There exists a collection of Morfisms, $\mathbf{Mor}(\mathbf{Q})$, denoted by $\mathbf{f}_{(\sigma)}, \mathbf{g}_{(\omega)}, \mathbf{h}_{(\tau)}, \dots$
- (3) For each morfism, $\mathbf{f}_{(\sigma)}$, there exists a *Domain* of $\mathbf{f}_{(\sigma)}$, $\mathbf{dom}(\mathbf{f}, \sigma) = \mathbf{A}$, and a *Codomain* of $\mathbf{f}_{(\sigma)}$, $\mathbf{cod}(\mathbf{f}, \sigma) = \mathbf{B}$, denoted by $\mathbf{f}_{(\sigma)}: \mathbf{A} \rightarrow \mathbf{B}$ or $\mathbf{A} \xrightarrow{\mathbf{f}_{(\sigma)}} \mathbf{B}$.
- (4) For each pair of morfisms, $\mathbf{f}_{(\sigma)}$ & $\mathbf{g}_{(\omega)}$, with $\mathbf{f}_{(\sigma)}: \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{g}_{(\omega)}: \mathbf{B} \rightarrow \mathbf{C}$ there exists a *Composition*, $\mathbf{g}_{(\omega)} \circ \mathbf{f}_{(\sigma)}$, such that $\mathbf{g}_{(\omega)} \circ \mathbf{f}_{(\sigma)}: \mathbf{A} \rightarrow \mathbf{C}$.
- (5) For each object, \mathbf{A} , there exists an *Identity*, $\mathbf{I}_{\mathbf{A}}$, such that $\mathbf{I}_{\mathbf{A}}: \mathbf{A} \rightarrow \mathbf{A}$.

RULES:

(a) ASSOCIATIVITY:

If $\mathbf{A} \xrightarrow{\mathbf{f}_{(\sigma)}} \mathbf{B} \xrightarrow{\mathbf{g}_{(\omega)}} \mathbf{C} \xrightarrow{\mathbf{h}_{(\tau)}} \mathbf{D}$, then $\mathbf{h}_{(\tau)} \circ (\mathbf{g}_{(\omega)} \circ \mathbf{f}_{(\sigma)}) = (\mathbf{h}_{(\tau)} \circ \mathbf{g}_{(\omega)}) \circ \mathbf{f}_{(\sigma)}$.

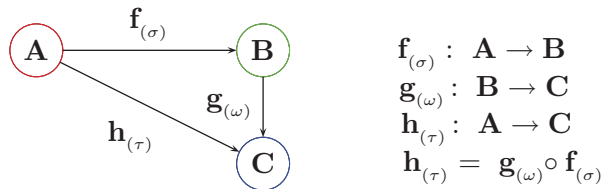
(b) IDENTITY:

If $\mathbf{A} \xrightarrow{\mathbf{f}_{(\sigma)}} \mathbf{B} \xrightarrow{\mathbf{g}_{(\omega)}} \mathbf{C}$, then $\mathbf{g}_{(\omega)} \circ \mathbf{I}_{\mathbf{B}} = \mathbf{g}_{(\omega)}$ and $\mathbf{I}_{\mathbf{B}} \circ \mathbf{f}_{(\sigma)} = \mathbf{f}_{(\sigma)}$.

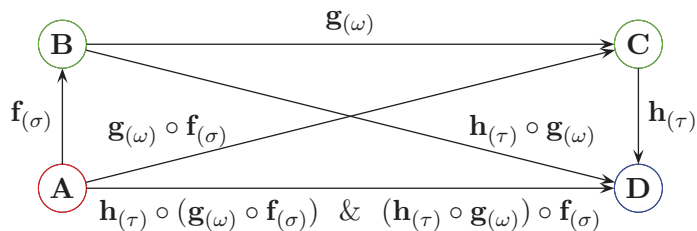
2.1. Kat Diagrams.


As with general Category theory, the formal properties relating Category objects and morfisms expressed above can also be expressed pictorially.

COMPOSITION: For $\mathbf{f}_{(\sigma)}: \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{g}_{(\omega)}: \mathbf{B} \rightarrow \mathbf{C}$ the composition, $\mathbf{g}_{(\omega)} \circ \mathbf{f}_{(\sigma)}: \mathbf{A} \rightarrow \mathbf{C}$, may be expressed pictorially by the following diagram.

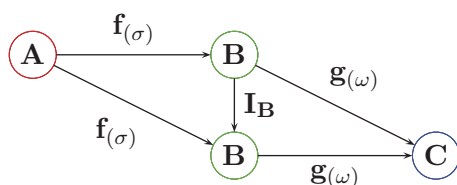


ASSOCIATIVITY: If $\mathbf{A} \xrightarrow{\mathbf{f}_{(\sigma)}} \mathbf{B} \xrightarrow{\mathbf{g}_{(\omega)}} \mathbf{C} \xrightarrow{\mathbf{h}_{(\tau)}} \mathbf{D}$, then $\mathbf{h}_{(\tau)} \circ (\mathbf{g}_{(\omega)} \circ \mathbf{f}_{(\sigma)}) = (\mathbf{h}_{(\tau)} \circ \mathbf{g}_{(\omega)}) \circ \mathbf{f}_{(\sigma)}$, may be expressed by the following diagram.



Definition 2.2. Identity: Every object, \mathbf{A} , has its own identity. $\mathbf{I}_A: \mathbf{A} \rightarrow \mathbf{A}$ 

IDENTITY: If $\mathbf{A} \xrightarrow{f(\sigma)} \mathbf{B} \xrightarrow{g(\omega)} \mathbf{C}$, then $g(\omega) \circ \mathbf{I}_B = g(\omega)$ and $\mathbf{I}_B \circ f(\sigma) = f(\sigma)$, may be expressed pictorially by the following diagram.



These diagrams mimic the familiar category diagrams.

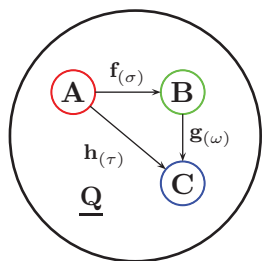
2.2. Kategories as Sets.

Kategories can be modeled by extended sets.

Definition 2.3. $Kat(\mathbf{Q})$: \mathbf{Q} is a Kategory iff:

- (a) $(\forall f, \sigma) (f \in_{\sigma} Mor(\mathbf{Q}) \ \& \ (\exists \mathbf{A}, \mathbf{B})(\mathbf{A} \in_{\mathbf{I}_A} Obj(\mathbf{Q}) \ \& \ \mathbf{B} \in_{\mathbf{I}_B} Obj(\mathbf{Q}))$
 $(\text{dom}(f, \sigma) = \mathbf{A}, \text{ and } \text{cod}(f, \sigma) = \mathbf{B})$,
- (b) $(\forall f, \sigma, g, \omega) (f \in_{\sigma} Mor(\mathbf{Q}) \ \& \ g \in_{\omega} Mor(\mathbf{Q}) \ \& \ \text{cod}(f, \sigma) = \text{dom}(g, \omega))$
 $\longrightarrow (\exists h, \tau) (h(\tau) = g(\omega) \circ f(\sigma) \ \& \ h \in_{\tau} Mor(\mathbf{Q}))$
- (c) $(\forall f, \sigma) (f \in_{\sigma} Mor(\mathbf{Q}) \ \& \ \text{dom}(f, \sigma) = \mathbf{A} \ \& \ \text{cod}(f, \sigma) = \mathbf{B})$
 $\longrightarrow (\mathbf{I}_B \circ f(\sigma) = f(\sigma) \circ \mathbf{I}_A \ \& \ \mathbf{A} \in_{\mathbf{I}_A} Obj(\mathbf{Q}) \ \& \ \mathbf{B} \in_{\mathbf{I}_B} Obj(\mathbf{Q}))$.

Example: $Kat(\mathbf{Q})$



\implies

$$\begin{aligned} \mathbf{Q} &= \{ \langle Obj(\mathbf{Q}), Mor(\mathbf{Q}) \rangle \} \\ Obj(\mathbf{Q}) &= \{ \mathbf{A}^{\mathbf{I}_A}, \mathbf{B}^{\mathbf{I}_B}, \mathbf{C}^{\mathbf{I}_C} \} \quad Mor(\mathbf{Q}) = \{ f^{\sigma}, g^{\omega}, h^{\tau} \} \\ \text{dom}(f, \sigma) &= \text{dom}(h, \tau) = \mathbf{A} \\ \text{cod}(f, \sigma) &= \text{dom}(g, \omega) = \mathbf{B} \\ \text{cod}(g, \omega) &= \text{cod}(h, \tau) = \mathbf{C} \\ h(\tau) &= g(\omega) \circ f(\sigma), \quad \tau = \langle \sigma_1, \omega_2 \rangle. \\ \mathbf{I}_C \circ h(\tau) \circ \mathbf{I}_A &= \mathbf{I}_C \circ g(\omega) \circ \mathbf{I}_B \circ f(\sigma) \circ \mathbf{I}_A \end{aligned}$$

Notice that by this construction $Obj(\mathbf{Q})$ is also the collection of *Identities* of $Kat(\mathbf{Q})$.